

Exact Real Computer Arithmetic Using Continued Fractions

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Motivation

$$\begin{pmatrix} 64919121 & -159018721 \\ 41869520.5 & -102558961 \end{pmatrix} x = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$

Haskell

```
solveAxB (a0, a1, a2, a3) (b0, b1) =  
  let det = a0 * a3 - a1 * a2  
  in ((a3 * b0 - a1 * b1) / det,  
      (-a2 * b0 + a0 * b1) / det)
```

Motivation

What Haskell (Ruby, Python, Java, ...) says...

$$x = \left(\begin{array}{c} 102558961 \\ 41869520.5 \end{array} \right)$$

The real answer...

$$x = \left(\begin{array}{c} 205117922 \\ 83739041 \end{array} \right)$$

Arbitrary Integer Arithmetic

- Many programming languages have arbitrary-precision **integer** arithmetic built-in.
- This can be used to implement arbitrary-precision floating point.
- However, it doesn't help us with $\sin(x)$, \sqrt{x} , etc.
- We want to compute with *exact* \mathbb{R} arithmetic.
 - ▶ And we want to be able to pick our own output precision and have the details handled for us.
- There are several solutions.

What is a *Continued Fraction*?

- Continued fractions represent a number (iteratively) as the sum of its integer part and the reciprocal of the remainder.
- For example: $\frac{5}{27} = 0 + \frac{1}{5 + \frac{1}{2 + \frac{1}{2 + 0}}}$
- We represent this by $[0; 5, 2, 2]$

A More Detailed Conversion Example

Converting 2.54 ($\frac{254}{100}$)

$$\frac{254}{100} = \mathbf{2} \text{ (remainder 54)}$$

$$\frac{100}{54} = \mathbf{1} \text{ (remainder 46)}$$

$$\frac{54}{46} = \mathbf{1} \text{ (remainder 8)}$$

$$\frac{46}{8} = \mathbf{5} \text{ (remainder 6)}$$

$$\frac{8}{6} = \mathbf{1} \text{ (remainder 2)}$$

$$\frac{6}{2} = \mathbf{3} \text{ (remainder 0)}$$

$$2.54 = \mathbf{2} + \frac{1}{\mathbf{1} + \frac{1}{\mathbf{1} + \frac{1}{\mathbf{5} + \frac{1}{\mathbf{1} + \frac{1}{\mathbf{3} + 0}}}}}$$

$$= [2; 1, 1, 5, 1, 3]$$

Some Properties Of Continued Fractions

- The continued fraction expansion of a number is finite if and only if the number is rational.
- Simple rational numbers have expansions with few (and small) terms.
- Truncating the expansion of a number gives a really good approximation of that number.

Some Properties Of Continued Fractions

as applied to programming

- Operations on them happen in “most-significant-digit first” order.
- Because of this, we can take full advantage of lazy evaluation in certain programming languages.
 - ▶ We can pick an accuracy level, then later decide we need more accuracy, and not have to recompute everything.

A More General Example

Let's expand:

$$\frac{70t + 29}{12t + 5}$$

... knowing only that t is positive.

We know the first term is 5, because as t varies between 0 and ∞ , $\frac{70}{12} < \frac{70t + 29}{12t + 5} < \frac{29}{5}$.

A More General Example

$$\frac{70t + 29}{12t + 5}$$

$$\frac{70t + 29}{12t + 5} = \mathbf{5} \text{ (remainder } 10t + 4)$$

$$\frac{12t + 5}{10t + 4} = \mathbf{1} \text{ (remainder } 2t + 1)$$

$$\frac{10t + 4}{2t + 1} = \mathbf{4} \text{ (remainder } 2t)$$

$$\begin{aligned} \frac{70t + 29}{12t + 5} &= \mathbf{5} + \frac{1}{\mathbf{1} + \frac{1}{\mathbf{4} + \frac{1}{\frac{2t+1}{2t}}}} \\ &= [5; 1, 4, \frac{2t+1}{2t}] \end{aligned}$$

...and we are stuck. All we know is that our next result is in $[1, \infty)$.

If we knew that $t > \frac{1}{2}$, then we could construct another term:

$$\frac{2t+1}{2t} = 1 + \frac{1}{2t}.$$

Going In Reverse

Input

Now suppose we are receiving terms of a continued fraction $([5; 1, 4, 1, \dots])$ which may go on forever or end with a symbolic expression.

Let's reconstruct the value as we receive terms.

Going In Reverse

Input

Let x be the “rest” of the continued fraction. Before we receive the first term, then, x represents the whole thing.

- When we get the first term (5), $x = 5 + \frac{1}{x'} = \frac{5x' + 1}{x'}$.
- When we get the next term (1), x becomes $1 + \frac{1}{x'}$ and $\frac{5x + 1}{x}$ becomes $\frac{6x' + 5}{x' + 1}$.
- When we get the next term (4), x becomes $4 + \frac{1}{x'}$ and $\frac{6x + 5}{x + 1}$ becomes $\frac{29x' + 6}{5x' + 1}$.
- When we get the next term (1), x becomes $1 + \frac{1}{x'}$ and $\frac{29x + 6}{5x + 1}$ becomes $\frac{35x' + 29}{6x' + 5}$.
- When x becomes $2t$, we have constructed the original value.

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Going In Reverse

The pattern

When x becomes $a + \frac{1}{x'}$, then $\frac{qx + r}{sx + t}$ becomes $\frac{(aq + r)x' + q}{(as + t)x' + s}$.

Arithmetic

Consider functions of the form $f(x) = \frac{ax + b}{cx + d}$ where $a, b, c, d \in \mathbb{Z}$.

These are called **homographic functions**.

We can express these in code as a 4-tuple.

Haskell

```
type CF = [Integer]
```

```
type Hom = (Integer, Integer,  
           Integer, Integer)
```

```
hom :: Hom -> CF -> CF    -- Goal.
```

Arithmetic

If $x = [x_0; \dots] = x_0 + \frac{1}{x'}$...
then by algebra,

$$\begin{aligned} \frac{\mathbf{ax} + \mathbf{b}}{\mathbf{cx} + \mathbf{d}} &= \frac{a(x_0 + \frac{1}{x'}) + b}{c(x_0 + \frac{1}{x'}) + d} \\ &= \frac{(\mathbf{ax}_0 + \mathbf{b})x' + \mathbf{a}}{(\mathbf{cx}_0 + \mathbf{d})x' + \mathbf{c}} \end{aligned}$$

Arithmetic

Absorbing terms

Absorbing terms

```

absorb :: Hom -> Integer -> Hom
absorb (a, b,
        c, d) x0 = (a * x0 + b, a,
                   c * x0 + d, c)

```

An invariant

```

hom h (x0:rest) == hom (absorb h x0) rest

```

Arithmetic

Emitting terms

If $z = \frac{ax + b}{cx + d}$, and $x \in \mathbb{Z}^+ \cup \{0\}$, then $z \in [\frac{a}{c}, \frac{b}{d}]$.

If $\frac{a}{c}$ and $\frac{b}{d}$ have the same integer part, q , then we know $z = [q; \dots]$.

Emitting terms

```
tryEmit :: Hom -> Maybe Integer
tryEmit (a, b,
        c, d) = if c /= 0 &&
                d /= 0 &&
                (a `quot` c == b `quot` d)
            then Just (a `quot` c)
            else Nothing
```

Arithmetic

Emitting terms

Let our result $z = q + \frac{1}{z'}$, then:

$$\begin{aligned}z' &= (z - q)^{-1} \\ &= \left(\frac{ax + b}{cx + d} - q \right)^{-1} \\ &= \frac{\mathbf{cx + d}}{\mathbf{(a - cq)x + (b - dq)}}\end{aligned}$$

Arithmetic

Emitting terms

Emitting terms

```
emit :: Hom -> Integer -> Hom
emit (a, b,
      c, d) q = (c,          d,
                 a - (c * q), b - (d * q))
```

An invariant

```
hom h cf == q : hom (emit q h) cf
```


Arithmetic

Deciding when to absorb or emit

Since we're using Haskell, we should take as much advantage of lazy evaluation as possible!

Finally, a `hom` implementation!

```
hom :: Hom -> CF -> CF
hom h cf@(x:rest) = case tryEmit h of
    Just z  -> z : hom (emit h z) cf
    Nothing -> hom (absorb h x) rest
```

Quick example

```
2 * pi == hom (2, 0,
              0, 1) pi
        == CF [6, 3, 1, 1, 7, 2, 146, 3, 6, 1, ...]
```

Back To Our Motivation

Before (using Double)

```
> solveAxB (64919121 :: Double, -159018721,  
           41869520.5,          -102558961) (1, 0)  
(1.02558961e8, 4.18695205e7)
```

After (using continued fractions)

```
> solveAxB (64919121 :: CF, -159018721,  
           41869520.5,          -102558961) (1, 0)  
(205117922, 83739041)
```

What I Haven't Shown

- Doing arithmetic on two continued fractions.
 - ▶ Same as we've done except for $\frac{axy + bx + cy + d}{exy + fx + gy + h}$.
 - ▶ These are called **bihomographic functions**.
 - ▶ By manipulating a through h , we can generate all of the arithmetic operations, and absorb/emit from x and y as before.
- Transcendental functions
- Integrating with Haskell's typeclass hierarchy

Future Research

- Implementation in Coq or another automated theorem assistant.
 - ▶ This would let us easily prove that our invariants hold in every possible case.
- Looking at continued fractions topologically.
- Employing type theory via domain-theoretic notions of definedness. (Peter Potts)

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